

From braids to transverse slices in reductive groups

Wicher Malten¹

¹Mathematical Institute
University of Oxford
Wicher.Malten@maths.ox.ac.uk

Lie Groups Seminar, Massachusetts Institute of Technology,
September 2021

Outline

- 1 Background
- 2 He-Lusztig's work
- 3 My work
- 4 End

One historical path

- Let G be a reductive group, fix a maximal torus T and denote the Weyl group by $W = N_G(T)/T$, and similarly for its Lie algebra \mathfrak{g} .

One historical path

- Let G be a reductive group, fix a maximal torus T and denote the Weyl group by $W = N_G(T)/T$, and similarly for its Lie algebra \mathfrak{g} .
- Around 1969, Grothendieck constructed a simultaneous resolution of the singularities of the fibres of

$$G \longrightarrow G//G \simeq T/W$$

and suggested that (strictly) transverse slices to conjugacy classes at subregular elements should yield universal deformations of the corresponding Du Val-Klein (or “ADE”) singularity, and similarly for the Lie algebra.

One historical path

- Let G be a reductive group, fix a maximal torus T and denote the Weyl group by $W = N_G(T)/T$, and similarly for its Lie algebra \mathfrak{g} .
- Around 1969, Grothendieck constructed a simultaneous resolution of the singularities of the fibres of

$$G \longrightarrow G//G \simeq T/W$$

and suggested that (strictly) transverse slices to conjugacy classes at subregular elements should yield universal deformations of the corresponding Du Val-Klein (or “ADE”) singularity, and similarly for the Lie algebra.

- Around 1980, Slowody constructed suitable slices in the Lie *algebra*, and studied this.

Slodowy slices

- They already appear in Harish-Chandra's 1964 work on invariant distributions on Lie algebras.

Slodowy slices

- They already appear in Harish-Chandra's 1964 work on invariant distributions on Lie algebras.
- They play a crucial role in the classification of certain infinite-dimensional representations appearing in the Langlands program ("Whittaker representations"), due to

Slodowy slices

- They already appear in Harish-Chandra's 1964 work on invariant distributions on Lie algebras.
- They play a crucial role in the classification of certain infinite-dimensional representations appearing in the Langlands program ("Whittaker representations"), due to
- The fact that they are the semi-classical limits of finite W -algebras (and their affine cousins).

Slodowy slices

- They already appear in Harish-Chandra's 1964 work on invariant distributions on Lie algebras.
- They play a crucial role in the classification of certain infinite-dimensional representations appearing in the Langlands program ("Whittaker representations"), due to
- The fact that they are the semi-classical limits of finite W -algebras (and their affine cousins).
- Have recently been applied to reconstruct Khovanov homology (Seidel-Smith, Abouzaid-Smith).

Slodowy slices

- They already appear in Harish-Chandra's 1964 work on invariant distributions on Lie algebras.
- They play a crucial role in the classification of certain infinite-dimensional representations appearing in the Langlands program ("Whittaker representations"), due to
- The fact that they are the semi-classical limits of finite W -algebras (and their affine cousins).
- Have recently been applied to reconstruct Khovanov homology (Seidel-Smith, Abouzaid-Smith).
- Appear in the work of numerous physicists on supersymmetric gauge theories (Gaiotto, Witten, etc.).

The Kostant Slice

- Kostant's slice: fixing a principal nilpotent element e in \mathfrak{g} , the Jacobson-Morozov theorem furnishes an embedding $\langle e, h, f \rangle = \mathfrak{sl}_2 \hookrightarrow \mathfrak{g}$; set $\mathfrak{s} := e + \ker \operatorname{ad} f \subset \mathfrak{g}$. It comes with two different cross section statements (from 1963 and 1978):

The Kostant Slice

- Kostant's slice: fixing a principal nilpotent element e in \mathfrak{g} , the Jacobson-Morozov theorem furnishes an embedding $\langle e, h, f \rangle = \mathfrak{sl}_2 \hookrightarrow \mathfrak{g}$; set $\mathfrak{s} := e + \ker \operatorname{ad} f \subset \mathfrak{g}$. It comes with two different cross section statements (from 1963 and 1978):
- The composition

$$\mathfrak{s} \hookrightarrow \mathfrak{g} \twoheadrightarrow \mathfrak{g} // G \simeq \mathfrak{t} / W$$

is an isomorphism.

The Kostant Slice

- Kostant's slice: fixing a principal nilpotent element e in \mathfrak{g} , the Jacobson-Morozov theorem furnishes an embedding $\langle e, h, f \rangle = \mathfrak{sl}_2 \hookrightarrow \mathfrak{g}$; set $\mathfrak{s} := e + \ker \operatorname{ad} f \subset \mathfrak{g}$. It comes with two different cross section statements (from 1963 and 1978):
- The composition

$$\mathfrak{s} \hookrightarrow \mathfrak{g} \twoheadrightarrow \mathfrak{g} // G \simeq \mathfrak{t} / W$$

is an isomorphism.

- Denote by $N_+ = [B, B]$ the unipotent radical of a Borel subgroup B , by N_- its opposite.

The Kostant Slice

- Kostant's slice: fixing a principal nilpotent element e in \mathfrak{g} , the Jacobson-Morozov theorem furnishes an embedding $\langle e, h, f \rangle = \mathfrak{sl}_2 \hookrightarrow \mathfrak{g}$; set $\mathfrak{s} := e + \ker \operatorname{ad} f \subset \mathfrak{g}$. It comes with two different cross section statements (from 1963 and 1978):
- The composition

$$\mathfrak{s} \hookrightarrow \mathfrak{g} \twoheadrightarrow \mathfrak{g} // G \simeq \mathfrak{t} / W$$

is an isomorphism.

- Denote by $N_+ = [B, B]$ the unipotent radical of a Borel subgroup B , by N_- its opposite.
- The adjoint action map

$$N_+ \times \mathfrak{s} \longrightarrow e + \mathfrak{n}_+^\perp =: \mu^{-1}(e)$$

is an isomorphism, where $\mathfrak{n}_+^\perp = \mathfrak{b}_+$ denotes the Killing form complement to \mathfrak{n}_+ .

The Steinberg Slice

- For $w \in W$, write

$$N_w := N_+ \cap w^{-1}N_-w = \prod_{\beta \in \mathfrak{R}_w} N_\beta,$$

where \mathfrak{R}_w is the set of positive roots made negative by w ,
and by T^w the points in T fixed by \dot{w} .

The Steinberg Slice

- For $w \in W$, write

$$N_w := N_+ \cap w^{-1}N_-w = \prod_{\beta \in \mathfrak{R}_w} N_\beta,$$

where \mathfrak{R}_w is the set of positive roots made negative by w ,
and by T^w the points in T fixed by \dot{w} .

- Steinberg's slice comes with similar cross sections (1965): if G is simply-connected and w a Coxeter element and $S := \dot{w}N_w$,

$$S \hookrightarrow G \twoheadrightarrow G//G \simeq T/W$$

is an isomorphism.

The Steinberg Slice

- For $w \in W$, write

$$N_w := N_+ \cap w^{-1}N_-w = \prod_{\beta \in \mathfrak{R}_w} N_\beta,$$

where \mathfrak{R}_w is the set of positive roots made negative by w , and by T^w the points in T fixed by \dot{w} .

- Steinberg's slice comes with similar cross sections (1965): if G is simply-connected and w a Coxeter element and $S := \dot{w}N_w$,

$$S \hookrightarrow G \twoheadrightarrow G//G \simeq T/W$$

is an isomorphism.

- Moreover, so is the conjugation action

$$N_+ \times S \longrightarrow N_+ \dot{w}N_+, \quad (n, s) \longmapsto n^{-1}sn$$

The Steinberg Slice

- For $w \in W$, write

$$N_w := N_+ \cap w^{-1}N_-w = \prod_{\beta \in \mathfrak{R}_w} N_\beta,$$

where \mathfrak{R}_w is the set of positive roots made negative by w , and by T^w the points in T fixed by \dot{w} .

- Steinberg's slice comes with similar cross sections (1965): if G is simply-connected and w a Coxeter element and $S := \dot{w}N_w$,

$$S \hookrightarrow G \twoheadrightarrow G//G \simeq T/W$$

is an isomorphism.

- Moreover, so is the conjugation action

$$N_+ \times S \longrightarrow N_+ \dot{w}N_+, \quad (n, s) \longmapsto n^{-1}sn$$

- (Proof of the second cross section is missing!!)

The Steinberg Slice

Example

Let $G = \mathrm{SL}_{r+1}$ over a commutative ring \mathcal{A} and consider the Coxeter element $w = s_1 \cdots s_r$. A suitable lift \dot{w} yields the Steinberg slice of Frobenius companion matrices

$$\dot{w}N_w = \left\{ \begin{bmatrix} 0 & 1 & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & 1 & 0 \\ 0 & 0 & \cdots & 0 & 1 \\ (-1)^r & c_r & \cdots & c_2 & c_1 \end{bmatrix} : c_1, \dots, c_r \in \mathcal{A} \right\}.$$

Generalisations

- Slodowy slices are constructed out of nilpotent elements, whereas Steinberg's slice is constructed out of Coxeter elements in the Weyl group.

Generalisations

- Slodowy slices are constructed out of nilpotent elements, whereas Steinberg's slice is constructed out of Coxeter elements in the Weyl group.
- In the late 1970s, Spaltenstein tried generalising Steinberg's slice to conjugates of Coxeter elements and noticed this second cross section property fails in type A_5 for

$$W = s_2 s_1 s_4 s_3 s_5 s_4 s_3 s_2 s_1.$$

Generalisations

- Slodowy slices are constructed out of nilpotent elements, whereas Steinberg's slice is constructed out of Coxeter elements in the Weyl group.
- In the late 1970s, Spaltenstein tried generalising Steinberg's slice to conjugates of Coxeter elements and noticed this second cross section property fails in type A_5 for

$$w = s_2 s_1 s_4 s_3 s_5 s_4 s_3 s_2 s_1.$$

- In 2011, Sevostyanov constructed slices out of Weyl group elements whose "eigenspaces" in the reflection representation can be ordered "nicely" w.r.t. the dominant Weyl chamber.

Generalisations

- Slodowy slices are constructed out of nilpotent elements, whereas Steinberg's slice is constructed out of Coxeter elements in the Weyl group.
- In the late 1970s, Spaltenstein tried generalising Steinberg's slice to conjugates of Coxeter elements and noticed this second cross section property fails in type A_5 for

$$w = s_2 s_1 s_4 s_3 s_5 s_4 s_3 s_2 s_1.$$

- In 2011, Sevostyanov constructed slices out of Weyl group elements whose "eigenspaces" in the reflection representation can be ordered "nicely" w.r.t. the dominant Weyl chamber.
- In 2012 (independently), He-Lusztig constructed slices of out *elliptic* Weyl group elements (= no fixed points in the reflection representation) which have minimal length.

An example

Example

Let $G = \mathrm{SL}_3$ over a commutative ring and $w := s_1 s_2 s_1$. The cross section statement asks whether the conjugation map

$$\left(\begin{pmatrix} 1 & n_1 & n_{12} \\ 0 & 1 & n_2 \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} x_1 & x_{12} & t \\ x_2 & -t^{-2} & 0 \\ t & 0 & 0 \end{pmatrix} \right) \in N_+ \times \dot{w} T^w N_+$$

to

$$\begin{pmatrix} n_{12}t + x_1 + n_1x_2 & -n_1t^{-2} + x_{12} - n_1(n_{12}t + x_1 + n_1x_2) & n_1n_2t^{-2} + t - n_2x_{12} + (n_1n_2 - n_{12})(n_{12}t + x_1 + n_1x_2) \\ n_2t + x_2 & -t^{-2} - n_1(n_2t + x_2) & n_2t^{-2} + (n_1n_2 - n_{12})(n_2t + x_2) \\ t & -n_1t & (n_1n_2 - n_{12})t \end{pmatrix}$$

in $N_+ \dot{w} T^w N_+$ is an isomorphism.

Dissimilarity

- Sevostyanov's conditions work for *some* elements in each conjugacy class

Dissimilarity

- Sevostyanov's conditions work for *some* elements in each conjugacy class

Lemma

But only for the two bipartite Coxeter elements

Dissimilarity

- Sevostyanov's conditions work for *some* elements in each conjugacy class

Lemma

But only for the two bipartite Coxeter elements

- He-Lusztig: all elliptic elements of minimal length, e.g.:

Dissimilarity

- Sevostyanov's conditions work for *some* elements in each conjugacy class

Lemma

But only for the two bipartite Coxeter elements

- He-Lusztig: all elliptic elements of minimal length, e.g.:

Example

All Coxeter elements are elliptic, and in type A all elliptic elements are conjugate to Coxeter elements.

Dissimilarity

- Sevostyanov's conditions work for *some* elements in each conjugacy class

Lemma

But only for the two bipartite Coxeter elements

- He-Lusztig: all elliptic elements of minimal length, e.g.:

Example

All Coxeter elements are elliptic, and in type A all elliptic elements are conjugate to Coxeter elements.

Outside of type A, there are always more.

Dissimilarity

- Sevostyanov's conditions work for *some* elements in each conjugacy class

Lemma

But only for the two bipartite Coxeter elements

- He-Lusztig: all elliptic elements of minimal length, e.g.:

Example

All Coxeter elements are elliptic, and in type A all elliptic elements are conjugate to Coxeter elements.

Outside of type A, there are always more.

- Sevostyanov's 2019 computations show that in order to construct strictly transverse slices to *all* conjugacy classes in reductive groups, you need to use most non-elliptic classes.

The braid monoid: definition

- Weyl groups are examples of finite Coxeter groups, which have a presentation

$$W = \langle s_1, \dots, s_{rk} : s_i s_j s_i \cdots = s_j s_i s_j \cdots, s_i^2 = 1 \rangle_{\text{grp}}$$

The braid monoid: definition

- Weyl groups are examples of finite Coxeter groups, which have a presentation

$$W = \langle s_1, \dots, s_{rk} : s_i s_j s_i \cdots = s_j s_i s_j \cdots, s_i^2 = 1 \rangle_{\text{grp}}$$

- The corresponding braid *monoid* is given by

$$B^+ := B_W^+ := \langle b_1, \dots, b_{rk} : b_i b_j b_i \cdots = b_j b_i b_j \cdots \rangle_{\text{mon}}$$

The braid monoid: definition

- Weyl groups are examples of finite Coxeter groups, which have a presentation

$$W = \langle s_1, \dots, s_{rk} : s_i s_j s_i \cdots = s_j s_i s_j \cdots, s_i^2 = 1 \rangle_{\text{grp}}$$

- The corresponding braid *monoid* is given by

$$B^+ := B_W^+ := \langle b_1, \dots, b_{rk} : b_i b_j b_i \cdots = b_j b_i b_j \cdots \rangle_{\text{mon}}$$

- The corresponding (Artin-Tits) braid *group* is given by

$$B := B_W := \langle b_1, \dots, b_{rk} : b_i b_j b_i \cdots = b_j b_i b_j \cdots \rangle_{\text{grp}}$$

The braid monoid: properties

- The braid monoid B^+ embeds into the braid group B .

The braid monoid: properties

- The braid monoid B^+ embeds into the braid group B .
- Moreover, any element in B can be expressed as a “fraction” of elements in B^+ .

The braid monoid: properties

- The braid monoid B^+ embeds into the braid group B .
- Moreover, any element in B can be expressed as a “fraction” of elements in B^+ .
- Matsumoto's theorem furnishes a well-defined inclusion of sets

$$W \longrightarrow B^+, \quad w \longmapsto b_w$$

by picking any *reduced* expression $w = s_{i_1} \cdots s_{i_l}$ and then mapping w to $b_{i_1} \cdots b_{i_l} =: b_{i_1 \cdots i_l} =: b_w$. The elements b_w are called *reduced/simple braids*.

The braid group: word problem

- Emil Artin (1925) wanted to construct for each element of B in type A a unique “word”, to be able to distinguish braids.

The braid group: word problem

- Emil Artin (1925) wanted to construct for each element of B in type A a unique “word”, to be able to distinguish braids.

Example

Let W be of type A_2 and consider

$$b_1 b_2 \stackrel{?}{=} b_2 b_1,$$

The braid group: word problem

- Emil Artin (1925) wanted to construct for each element of B in type A a unique “word”, to be able to distinguish braids.

Example

Let W be of type A_2 and consider

$$b_1 b_2 \stackrel{?}{=} b_2 b_1,$$

$$b_1 b_2 b_1 \stackrel{?}{=} b_2 b_1 b_2.$$

The braid group: example

Example

Let W be of type A_4 and consider

The braid group: example

Example

Let W be of type A_4 and consider

$$(b_1 b_2 b_1 b_3 b_2 b_4)^3 \stackrel{?}{=} b_1 b_2 b_3 b_4 b_1 b_2 b_3 b_1 b_2 b_1 b_3 b_4 b_2 b_2 b_3 b_4 b_1 b_2.$$

Deligne-Garside normal form: definition

- Artin found a solution in type A (“braid/Artin combing”)

Deligne-Garside normal form: definition

- Artin found a solution in type A (“braid/Artin combing”)
- He then wanted to know when elements are conjugate, but this “word” did not tell him when braids are conjugate

Deligne-Garside normal form: definition

- Artin found a solution in type A (“braid/Artin combing”)
- He then wanted to know when elements are conjugate, but this “word” did not tell him when braids are conjugate
- Garside gave a new solution to the word problem that *also* solves the conjugacy problem (~1965).

Deligne-Garside normal form: definition

- Artin found a solution in type A (“braid/Artin combing”)
- He then wanted to know when elements are conjugate, but this “word” did not tell him when braids are conjugate
- Garside gave a new solution to the word problem that *also* solves the conjugacy problem (~1965).
- Roughly speaking, the (right) Deligne-Garside normal form of a b braid in B^+ is obtained by decomposing it as a product of reduced braids $b = b_{w_n} \cdots b_{w_1}$, and then making the rightmost factors as large as possible.

Deligne-Garside normal form: definition

- Artin found a solution in type A (“braid/Artin combing”)
- He then wanted to know when elements are conjugate, but this “word” did not tell him when braids are conjugate
- Garside gave a new solution to the word problem that *also* solves the conjugacy problem (~1965).
- Roughly speaking, the (right) Deligne-Garside normal form of a b braid in B^+ is obtained by decomposing it as a product of reduced braids $b = b_{w_n} \cdots b_{w_1}$, and then making the rightmost factors as large as possible.
- So apparently this yields a unique expression

$$b_{w_m} \cdots b_{w_1} =: \text{DG}_m(b) \cdots \text{DG}_1(b).$$

Deligne-Garside normal form: definition

- Artin found a solution in type A (“braid/Artin combing”)
- He then wanted to know when elements are conjugate, but this “word” did not tell him when braids are conjugate
- Garside gave a new solution to the word problem that *also* solves the conjugacy problem (~1965).
- Roughly speaking, the (right) Deligne-Garside normal form of a b braid in B^+ is obtained by decomposing it as a product of reduced braids $b = b_{w_n} \cdots b_{w_1}$, and then making the rightmost factors as large as possible.
- So apparently this yields a unique expression

$$b_{w_m} \cdots b_{w_1} =: \text{DG}_m(b) \cdots \text{DG}_1(b).$$

- We write $\text{DG}(b) := \text{DG}_1(b)$, and will often identify it with the corresponding Coxeter group element w_1 .

Deligne-Garside normal form: back to examples

Example

Let W be of type A_2 , now find

$$b_1 b_2 = b_{s_1} b_{s_2} = b_{s_1 s_2} \neq b_{s_2 s_1} = b_2 b_1,$$

Deligne-Garside normal form: back to examples

Example

Let W be of type A_2 , now find

$$b_1 b_2 = b_{s_1} b_{s_2} = b_{s_1 s_2} \neq b_{s_2 s_1} = b_2 b_1,$$

$$b_2 b_1 b_2 = b_{212} = b_{121} = b_1 b_2 b_1.$$

Deligne-Garside normal form: back to examples

Example

Let W be of type A_2 , now find

$$b_1 b_2 = b_{s_1} b_{s_2} = b_{s_1 s_2} \neq b_{s_2 s_1} = b_2 b_1,$$

$$b_2 b_1 b_2 = b_{212} = b_{121} = b_1 b_2 b_1.$$

Example

Let W be of type A_4 , now find

$$\begin{aligned} (b_1 b_2 b_1 b_3 b_2 b_4)^3 &= b_{23} b_{341231} b_{w_0} \\ &= b_1 b_2 b_3 b_4 b_1 b_2 b_3 b_1 b_2 b_1 b_3 b_4 b_2 b_2 b_3 b_4 b_1 b_2. \end{aligned}$$

He-Lusztig's result

- Recall: Steinberg's claim is for Coxeter elements, e.g. $s_1 \cdots s_{rk}$ where rk is the rank of W (or G): the conjugation action

$$N_+ \times \dot{w}N_w \xrightarrow{\sim} N_+ \dot{w}N_+$$

is an isomorphism.

He-Lusztig's result

- Recall: Steinberg's claim is for Coxeter elements, e.g. $s_1 \cdots s_{rk}$ where rk is the rank of W (or G): the conjugation action

$$N_+ \times \dot{w}N_w \xrightarrow{\sim} N_+ \dot{w}N_+$$

is an isomorphism.

- The cross sections of He-Lusztig apply to elliptic elements w of minimal length in their conjugacy class, in the same way:

$$N_+ \times \dot{w}N_w \xrightarrow{\sim} N_+ \dot{w}N_+.$$

He-Lusztig's proof

- Ultimately, consists of two major steps:

He-Lusztig's proof

- Ultimately, consists of two major steps:
- (1): Proven “directly” for all elements w , such that $DG(b_w^d) = w_o$ for some integer $d \geq 1$. From case-by-case work (Geck-Michel), it was known then that this is true for *some* elements of minimal length in each elliptic conjugacy class, when $d = \text{ord}(w)$.

He-Lusztig's proof

- Ultimately, consists of two major steps:
- (1): Proven “directly” for all elements w , such that $DG(b_w^d) = w_o$ for some integer $d \geq 1$. From case-by-case work (Geck-Michel), it was known then that this is true for *some* elements of minimal length in each elliptic conjugacy class, when $d = \text{ord}(w)$.
- (2): If it is true for an element $w = xy$ with $\ell(w) = \ell(x) + \ell(y)$, then it is also true for $w' := yx$ if $\ell(y) + \ell(x) = \ell(w')$. From case-by-case work (Geck-Pfeiffer), it was known then that all elliptic elements of minimal length are conjugate to each other by such *cyclic shifts*.

He-Lusztig's proof

- Ultimately, consists of two major steps:
- (1): Proven “directly” for all elements w , such that $\text{DG}(b_w^d) = w_\circ$ for some integer $d \geq 1$. From case-by-case work (Geck-Michel), it was known then that this is true for *some* elements of minimal length in each elliptic conjugacy class, when $d = \text{ord}(w)$.
- (2): If it is true for an element $w = xy$ with $\ell(w) = \ell(x) + \ell(y)$, then it is also true for $w' := yx$ if $\ell(y) + \ell(x) = \ell(w')$. From case-by-case work (Geck-Pfeiffer), it was known then that all elliptic elements of minimal length are conjugate to each other by such *cyclic shifts*.
- (2'): Simpler: if w and w' are conjugate by cyclic shifts and $\text{DG}(b_w^d) = w_\circ$, then $\text{DG}(b_{w'}^{d'}) = w_\circ$ for some d' .

He-Lusztig's proof, tweaked

- So it becomes, over any ring:

He-Lusztig's proof, tweaked

- So it becomes, over any ring:
- (1): Cross section holds if $\text{DG}(b_w^d) = w_\circ$ for some computable integer $d \geq 1$, say $d = |\mathfrak{R}_+| - \ell(w) + 1$.

He-Lusztig's proof, tweaked

- So it becomes, over any ring:
- (1): Cross section holds if $DG(b_w^d) = w_o$ for some computable integer $d \geq 1$, say $d = |\mathfrak{R}_+| - \ell(w) + 1$.
- (2): This braid equation holds for all elliptic elements of minimal length.

He-Lusztig's proof, tweaked

- So it becomes, over any ring:
- (1): Cross section holds if $\text{DG}(b_w^d) = w_o$ for some computable integer $d \geq 1$, say $d = |\mathfrak{R}_+| - \ell(w) + 1$.
- (2): This braid equation holds for all elliptic elements of minimal length.

Lemma

Sevostyanov's elliptic elements satisfy this braid equation.

He-Lusztig's proof, tweaked

- So it becomes, over any ring:
- (1): Cross section holds if $\mathrm{DG}(b_w^d) = w_\circ$ for some computable integer $d \geq 1$, say $d = |\mathfrak{R}_+| - \ell(w) + 1$.
- (2): This braid equation holds for all elliptic elements of minimal length.

Lemma

Sevostyanov's elliptic elements satisfy this braid equation.

- Do his non-elliptic satisfy it? Rarely... but those slices are a bit different!

New definitions: firmly convex elements

- Let W be a finite Coxeter group. An element w is called *firmly convex* if:

New definitions: firmly convex elements

- Let W be a finite Coxeter group. An element w is called *firmly convex* if:
 - the subset of roots \mathfrak{R}^w that it fixes, forms a standard parabolic subroot system.

New definitions: firmly convex elements

- Let W be a finite Coxeter group. An element w is called *firmly convex* if:
 - the subset of roots \mathfrak{R}^w that it fixes, forms a standard parabolic subroot system.
 - + technical condition.

New definitions: firmly convex elements

- Let W be a finite Coxeter group. An element w is called *firmly convex* if:
 - the subset of roots \mathfrak{R}^w that it fixes, forms a standard parabolic subroot system.
 - + technical condition.

Lemma

\mathfrak{R}^w forms a standard parabolic subsystem if and only if the complement $\mathfrak{R}_+ \setminus \mathfrak{R}^w$ is convex, i.e.:

If $\beta_0, \beta_1 \in \mathfrak{R}_+ \setminus \mathfrak{R}^w$ and $c_0, c_1 \in \mathbb{R}_{>0}$ are such that $c_0\beta_0 + c_1\beta_1$ is again a root, then it lies in $\mathfrak{R}_+ \setminus \mathfrak{R}^w$.

New definitions: braid power bound

Definition

Let w_o denote the longest element of W . Given a firmly convex element w , let w_f denote the longest element of the standard parabolic subsystem \mathfrak{A}^w ; this yields a *braid power bound*

$$w_o w_f.$$

New definitions: braid power bound

Definition

Let w_o denote the longest element of W . Given a firmly convex element w , let w_f denote the longest element of the standard parabolic subsystem \mathfrak{R}^w ; this yields a *braid power bound*

$$w_o w_f.$$

Example

Let W be of type A_3 . If w is reflecting in $\alpha_1 + \alpha_2 + \alpha_3$, then this is $w_o s_2$.

New definitions: braid power bound

Definition

Let w_o denote the longest element of W . Given a firmly convex element w , let w_f denote the longest element of the standard parabolic subsystem \mathfrak{R}^w ; this yields a *braid power bound*

$$w_o w_f.$$

Example

Let W be of type A_3 . If w is reflecting in $\alpha_1 + \alpha_2 + \alpha_3$, then this is $w_o s_2$.

- So $\mathfrak{R}_{w_o w_f} = \mathfrak{R}_+ \setminus \mathfrak{R}^w$.

New definitions: dominant elements

Definition

Let C denote the dominant Weyl chamber. For any w , let $V_w = \text{im}(\text{id} - w)$ denote the orthogonal complement to the subset of fixed points $\ker(\text{id} - w)$ in the reflection representation.

New definitions: dominant elements

Definition

Let C denote the dominant Weyl chamber. For any w , let $V_w = \text{im}(\text{id} - w)$ denote the orthogonal complement to the subset of fixed points $\ker(\text{id} - w)$ in the reflection representation. Then w is called *dominant* if the closure \overline{C} of C contains an open subset of V_w .

New definitions: dominant elements

Definition

Let C denote the dominant Weyl chamber. For any w , let $V_w = \text{im}(\text{id} - w)$ denote the orthogonal complement to the subset of fixed points $\ker(\text{id} - w)$ in the reflection representation. Then w is called *dominant* if the closure \overline{C} of C contains an open subset of V_w .

Example

Reflection in a root is dominant if and only if this root is the highest root or the highest short root.

New definitions: dominant elements

Lemma

An involution has maximal length if and only if it is dominant.

New definitions: dominant elements

Lemma

An involution has maximal length if and only if it is dominant.

Lemma

For any element w there are implications

elliptic or Sevostyanov element \implies dominant \implies firmly convex

Transversality

- Let G be a manifold (or variety), and let C and S be two submanifolds. We say that the intersection $C \cap S$ is *transverse* if for all $g \in C \cap S$, we have

$$T_g G = T_g C + T_g S.$$

Transversality

- Let G be a manifold (or variety), and let C and S be two submanifolds. We say that the intersection $C \cap S$ is *transverse* if for all $g \in C \cap S$, we have

$$T_g G = T_g C + T_g S.$$

- We say that the intersection is *strictly transverse* if this is a direct sum, i.e.

$$T_g C \cap T_g S = \{0\}.$$

Inspiration from braids

- For any w , analysing roots shows that

$$\mathfrak{A}_{\text{DG}(b_w^d)} \subseteq \mathfrak{A}_+ \setminus \mathfrak{A}^w.$$

Inspiration from braids

- For any w , analysing roots shows that

$$\mathfrak{R}_{\text{DG}(b_w^d)} \subseteq \mathfrak{R}_+ \setminus \mathfrak{R}^w.$$

- By the “convexity” lemma, this inclusion is strict if w is not firmly convex; if it is firmly convex then it is equivalent to

$$\text{DG}(b_w^d) \leq w_o w_f$$

in the left weak Bruhat-Chevalley order.

Inspiration from braids

- For any w , analysing roots shows that

$$\mathfrak{R}_{\text{DG}(b_w^d)} \subseteq \mathfrak{R}_+ \setminus \mathfrak{R}^w.$$

- By the “convexity” lemma, this inclusion is strict if w is not firmly convex; if it is firmly convex then it is equivalent to

$$\text{DG}(b_w^d) \leq w_o w_f$$

in the left weak Bruhat-Chevalley order.

- We can modify Sevostyanov's definitions to come up with a cross section *statement*

$$N \times \dot{w}L^w N_w \longrightarrow N \dot{w}L^w N,$$

for any firmly convex element w . Here $N \subseteq N_+$ is generated by root subgroups for roots in $\mathfrak{R}_+ \setminus \mathfrak{R}^w$, whereas L^w is the reductive subgroup “generated” by \mathfrak{R}^w and T^w .

From braids to cross sections

- Can now modify the He-Lusztig proof to:

From braids to cross sections

- Can now modify the He-Lusztig proof to:

Theorem

If w is firmly convex and for some $d \geq 1$ we have

$$\mathrm{DG}(b_w^d) = w_o w_f,$$

then the conjugation map

$$N \times \dot{w}L^w N_w \longrightarrow N \dot{w}L^w N, \quad (n, s) \longmapsto n^{-1}sn$$

is an isomorphism, over any commutative ring.

From braids to cross sections

- Can now modify the He-Lusztig proof to:

Theorem

If w is firmly convex and for some $d \geq 1$ we have

$$\mathrm{DG}(b_w^d) = w_o w_f,$$

then the conjugation map

$$N \times \dot{w}L^w N_w \longrightarrow N \dot{w}L^w N, \quad (n, s) \longmapsto n^{-1}sn$$

is an isomorphism, over any commutative ring.

Lemma

He-Lusztig's and Sevostyanov's elements satisfy this equation.

More?

- How about Poisson structures?

More?

- How about Poisson structures?
- How about transversality?

More?

- How about Poisson structures?
- How about transversality?
- How about strict transversality?

More?

- How about Poisson structures?
- How about transversality?
- How about strict transversality?
- How about the converse?

Poisson structures

- In the late 90s, some people tried to obtain quantum analogues of finite W -algebras, but failed to obtain suitable characters.

Poisson structures

- In the late 90s, some people tried to obtain quantum analogues of finite W -algebras, but failed to obtain suitable characters.
- Sevostyanov '99 succeeded, by slightly “modifying” the Drinfeld-Jimbo quantum group $U_q G$.

Poisson structures

- In the late 90s, some people tried to obtain quantum analogues of finite W -algebras, but failed to obtain suitable characters.
- Sevostyanov '99 succeeded, by slightly “modifying” the Drinfeld-Jimbo quantum group $U_q G$.
- Can reinterpret his solution as a Drinfeld twist.

Poisson structures

- In the late 90s, some people tried to obtain quantum analogues of finite W -algebras, but failed to obtain suitable characters.
- Sevostyanov '99 succeeded, by slightly “modifying” the Drinfeld-Jimbo quantum group $U_q G$.
- Can reinterpret his solution as a Drinfeld twist.
- Quasiclassically, this twist corresponds to modifying the Semenov-Tian-Shansky bracket on G . Using the cross section isomorphism, can show:

Poisson structures

- In the late 90s, some people tried to obtain quantum analogues of finite W -algebras, but failed to obtain suitable characters.
- Sevostyanov '99 succeeded, by slightly “modifying” the Drinfeld-Jimbo quantum group $U_q G$.
- Can reinterpret his solution as a Drinfeld twist.
- Quasiclassically, this twist corresponds to modifying the Semenov-Tian-Shansky bracket on G . Using the cross section isomorphism, can show:

Lemma

This Poisson bracket reduces to a Poisson bracket on the slices if and only if such a twist is made.

Transversality again

- Sevostyanov deduces transversality by combining the cross section statement for w and the cross section statement for w^{-1} .

Transversality again

- Sevostyanov deduces transversality by combining the cross section statement for w and the cross section statement for w^{-1} .
- So we would need: $DG(b_w^d) = w_o w_f$ if and only if $DG(b_{w^{-1}}^{d'}) = w_o w_f$.

Transversality again

- Sevostyanov deduces transversality by combining the cross section statement for w and the cross section statement for w^{-1} .
- So we would need: $DG(b_w^d) = w_o w_f$ if and only if $DG(b_{w^{-1}}^{d'}) = w_o w_f$.

Example

Consider $w = s_3 s_1 s_2 s_3$ in type B_3 ; it does not fix any roots so it is closed, but for any integer $d > 1$ we have

$$DGN(b_w^d) = b_w^d \quad \text{and} \quad DGN(b_{w^{-1}}^d) = b_{323} b_w^{d-2} b_{13213}.$$

Transversality again

- Sevostyanov deduces transversality by combining the cross section statement for w and the cross section statement for w^{-1} .
- So we would need: $DG(b_w^d) = w_o w_f$ if and only if $DG(b_{w^{-1}}^{d'}) = w_o w_f$.

Example

Consider $w = s_3 s_1 s_2 s_3$ in type B_3 ; it does not fix any roots so it is closed, but for any integer $d > 1$ we have

$$DGN(b_w^d) = b_w^d \quad \text{and} \quad DGN(b_{w^{-1}}^d) = b_{323} b_w^{d-2} b_{13213}.$$

- We will see that this is true, with $d' = d$. Surprising... because normally $DG(b_w^d)$ and $DG(b_{w^{-1}}^d)$ are very different!

The converse

- The cross section statement is almost a statement about roots.

The converse

- The cross section statement is almost a statement about roots.
- But what is the identity $DG(b_w^d) = w_o w_f$ *really* doing in the proof?

The converse

- The cross section statement is almost a statement about roots.
- But what is the identity $DG(b_w^d) = w_o w_f$ really doing in the proof?
- It's trying to make all the roots in $\mathfrak{R}_+ \setminus \mathfrak{R}^w$ negative, step by step:

$$\begin{aligned} DG(b_w^d) = w_o w_f &\implies \text{cross}_w^d(\mathfrak{R}_+ \setminus \mathfrak{R}^w) = \emptyset \\ &\implies \text{cross section is isomorphism} \end{aligned}$$

Crossing roots

Definition

For any positive root β and w , we obtain a subset of positive roots

$$\text{cross}_w(\beta) := \left\{ w\left(\beta + \sum_{i=1}^m \beta_i\right) \in \mathfrak{R} : \beta_1, \dots, \beta_m \in \mathfrak{R}_w, m \geq 0 \right\} \cap \mathfrak{R}_+$$

and for a subset of positive roots $\mathfrak{N} \subseteq \mathfrak{R}_+$ we set

$$\text{cross}_w(\mathfrak{N}) := \bigcup_{\beta \in \mathfrak{N}} \text{cross}_w(\beta).$$

Crossing roots

Definition

For any positive root β and w , we obtain a subset of positive roots

$$\text{cross}_w(\beta) := \left\{ w\left(\beta + \sum_{i=1}^m \beta_i\right) \in \mathfrak{R} : \beta_1, \dots, \beta_m \in \mathfrak{R}_w, m \geq 0 \right\} \cap \mathfrak{R}_+$$

and for a subset of positive roots $\mathfrak{N} \subseteq \mathfrak{R}_+$ we set

$$\text{cross}_w(\mathfrak{N}) := \bigcup_{\beta \in \mathfrak{N}} \text{cross}_w(\beta).$$

Example

What is $\text{cross}_w(\beta)$ when β lies in \mathfrak{R}_w ? When $w(\beta)$ is simple?

Crossing roots

Lemma

- *For any simple root α not in \mathfrak{R}_w , the set $\text{cross}_w(\alpha)$ contains simple roots.*

Crossing roots

Lemma

- For any simple root α not in \mathfrak{R}_w , the set $\text{cross}_w(\alpha)$ contains simple roots.
- Implies: For any other element v of W and integer $d \geq 0$,

$$\text{DG}(b_w^d) \geq v \quad \text{if and only if} \quad \text{cross}_w^d(\mathfrak{R}_v) = \emptyset,$$

if and only if $\text{cross}_w^d(\mathfrak{R}_v)$ does not contain any simple roots.

Crossing roots

Lemma

- For any simple root α not in \mathfrak{R}_w , the set $\text{cross}_w(\alpha)$ contains simple roots.
- Implies: For any other element v of W and integer $d \geq 0$,

$$\text{DG}(b_w^d) \geq v \quad \text{if and only if} \quad \text{cross}_w^d(\mathfrak{R}_v) = \emptyset,$$

if and only if $\text{cross}_w^d(\mathfrak{R}_v)$ does not contain any simple roots.

- In particular: w is firmly convex and satisfies the braid equation $\text{DG}(b_w^d) = w \circ w_f$ if and only if $\text{cross}_w^d(\mathfrak{R}_+ \setminus \mathfrak{R}^w) = \emptyset$.

Crossing roots

Lemma

- For any simple root α not in \mathfrak{R}_w , the set $\text{cross}_w(\alpha)$ contains simple roots.
- Implies: For any other element v of W and integer $d \geq 0$,

$$\text{DG}(b_w^d) \geq v \quad \text{if and only if} \quad \text{cross}_w^d(\mathfrak{R}_v) = \emptyset,$$

if and only if $\text{cross}_w^d(\mathfrak{R}_v)$ does not contain any simple roots.

- In particular: w is firmly convex and satisfies the braid equation $\text{DG}(b_w^d) = w_0 w_f$ if and only if $\text{cross}_w^d(\mathfrak{R}_+ \setminus \mathfrak{R}^w) = \emptyset$.
- And that easily implies: if w is firmly convex then $\text{DG}(b_w^d) = w_0 w_f$ if and only if $\text{DG}(b_{w^{-1}}^d) = w_0 w_f$.

Strict transversality: minimally dominant elements

Definitions

A dominant element is called *minimally dominant* if its length is minimal among the dominant elements in its conjugacy class.

Strict transversality: minimally dominant elements

Definitions

A dominant element is called *minimally dominant* if its length is minimal among the dominant elements in its conjugacy class.

Example

For elliptic conjugacy classes, “minimally dominant” = “has minimal length”.

Strict transversality: minimally dominant elements

Definitions

A dominant element is called *minimally dominant* if its length is minimal among the dominant elements in its conjugacy class.

Example

For elliptic conjugacy classes, “minimally dominant” = “has minimal length”.

Lemma

For (nontrivial) non-elliptic conjugacy classes, minimally dominant elements never have minimal length.

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$
- Geck-Pfeiffer/He-Nie: Elliptic elements of minimal length are conjugate by cyclic shifts

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$
- Geck-Pfeiffer/He-Nie: Elliptic elements of minimal length are conjugate by cyclic shifts
- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ$ for some d

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$
- Geck-Pfeiffer/He-Nie: Elliptic elements of minimal length are conjugate by cyclic shifts
- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ$ for some d

Lemma

Every conjugacy class contains a minimally dominant element w such that $DG(b_w^{\text{ord}(w)}) = w_\circ w_f$, and

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$
- Geck-Pfeiffer/He-Nie: Elliptic elements of minimal length are conjugate by cyclic shifts
- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ$ for some d

Lemma

Every conjugacy class contains a minimally dominant element w such that $DG(b_w^{\text{ord}(w)}) = w_\circ w_f$, and minimally dominant elements are conjugate by cyclic shifts

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$
- Geck-Pfeiffer/He-Nie: Elliptic elements of minimal length are conjugate by cyclic shifts
- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ$ for some d

Lemma

Every conjugacy class contains a minimally dominant element w such that $DG(b_w^{\text{ord}(w)}) = w_\circ w_f$, and minimally dominant elements are conjugate by cyclic shifts

- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ w_f$ for some d

Braid powers of minimally dominant elements

- Geck-Michel/He-Nie: Every elliptic conjugacy class contains an element w of minimal length such that $DG(b_w^{\text{ord}(w)}) = w_\circ$
- Geck-Pfeiffer/He-Nie: Elliptic elements of minimal length are conjugate by cyclic shifts
- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ$ for some d

Lemma

Every conjugacy class contains a minimally dominant element w such that $DG(b_w^{\text{ord}(w)}) = w_\circ w_f$, and minimally dominant elements are conjugate by cyclic shifts

- Combine: \Rightarrow they all satisfy $DG(b_w^d) = w_\circ w_f$ for some d
- So by the previous theorem, they all yield transverse slices!

Strict transversality and minimally dominant elements

- In 2019, Sevostyanov showed that a subset of his elements yield strictly transverse slices (by using Lusztig's partition).

Strict transversality and minimally dominant elements

- In 2019, Sevostyanov showed that a subset of his elements yield strictly transverse slices (by using Lusztig's partition).
- He already knew they were transverse, so his main ingredient is a case-by-case dimension calculation.

Strict transversality and minimally dominant elements

- In 2019, Sevostyanov showed that a subset of his elements yield strictly transverse slices (by using Lusztig's partition).
- He already knew they were transverse, so his main ingredient is a case-by-case dimension calculation.
- Can show that these elements are all minimally dominant.

Strict transversality and minimally dominant elements

- In 2019, Sevostyanov showed that a subset of his elements yield strictly transverse slices (by using Lusztig's partition).
- He already knew they were transverse, so his main ingredient is a case-by-case dimension calculation.
- Can show that these elements are all minimally dominant.
- Can now deduce that all minimally dominant elements in these conjugacy classes yield strictly transverse slices!

Final statement

Theorem

Let C be a conjugacy class of a connected reductive group over an algebraically closed field, and let w be a minimally dominant element in the corresponding conjugacy class in Lusztig's partition.

Final statement

Theorem

Let C be a conjugacy class of a connected reductive group over an algebraically closed field, and let w be a minimally dominant element in the corresponding conjugacy class in Lusztig's partition.

Then C is strictly transversally intersected by $wL^w N_w$, and this slice inherits a natural Poisson structure.

End

- Thanks for listening!!
- Questions? Ideas??
- w.malten@gmail.com